

Positive k -spanning sets and their use in Derivative Free Optimization

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- 1 Positive Spanning Sets
- 2 Cosine measure
- 3 Positive k-Spanning Sets
- 4 k-cosine measure

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2 Cosine measure

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Applications

Question

Let $f : \mathbb{R}^n \rightarrow \mathbb{R}$ be a smooth function. How can we minimize $f(x)$?

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If the elements of D are well spread in \mathbb{R}^n , the algorithm will converge.

Positive Spanning Sets and Positive Bases

Definition

A **PSS** spans \mathbb{R}^n with positive linear combinations.

A **positive basis** is minimal for this property.

Remark

If the columns of matrix M form a PSS, we say that M is a PSS.

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Examples in \mathbb{R}^2

$\begin{bmatrix} 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{bmatrix}$ is a *positive basis*.

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$\begin{bmatrix} 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{bmatrix}$ is a *positive basis*. $\begin{bmatrix} 1 & 0 & -1 \\ 0 & 1 & -1 \end{bmatrix}$ *too*.

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Positive Spanning Sets and Positive Bases

Definition

A **PSS** spans \mathbb{R}^n with positive linear combinations. (or non-negative.)
A **positive basis** is minimal for this property.

Examples in \mathbb{R}^2

$\begin{bmatrix} 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{bmatrix}$ is a positive basis. $\begin{bmatrix} 1 & 0 & -1 \\ 0 & 1 & -1 \end{bmatrix}$ too.

e.g: $\begin{bmatrix} 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{bmatrix} \begin{bmatrix} 3 \\ 0 \\ 0 \\ 4 \end{bmatrix} = \begin{bmatrix} 3 \\ -4 \end{bmatrix}$ and $\begin{bmatrix} 1 & 0 & -1 \\ 0 & 1 & -1 \end{bmatrix} \begin{bmatrix} 7 \\ 0 \\ 4 \end{bmatrix} = \begin{bmatrix} 3 \\ -4 \end{bmatrix}$.

Remark

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$\begin{bmatrix} 1 & 0 & -2 & 1 \\ 0 & 1 & 1 & -2 \end{bmatrix}$ *is...*

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Remark

If the columns of matrix M form a PSS, we say that M is a PSS.

Cardinal

Proposition

*The smallest possible size for a positive basis of \mathbb{R}^n is $n + 1$.
The biggest possible size is $2n$.*

Proof (idea)

- *Lower bound: a PSS must clearly be larger than a linear basis.*
- *Upper bound: Trickier. Proven using linear programming arguments.*

Farkas Lemma

Lemma

Let $M \in \mathbb{R}^{n,m}$, let $b \in \mathbb{R}^n$. Exactly one of the two following assertions is true:

- Equation $Mx = b$ has a solution $x \geq 0$.
- Inequation $y^T M \geq 0$ has a solution y such that $y^T b < 0$.

Farkas Lemma

Lemma

Let $M \in \mathbb{R}^{n,m}$, let $b \in \mathbb{R}^n$. Exactly one of the two following assertions is true:

- Equation $Mx = b$ has a solution $x \geq 0$. *Always true for PSSs.*
- Inequation $y^T M \geq 0$ has a solution y such that $y^T b < 0$.

Remark

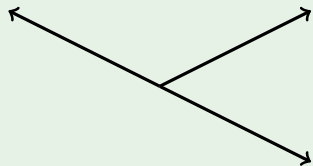
When M is a PSS, the second assertion is false for all b !

Characterization of PSSs

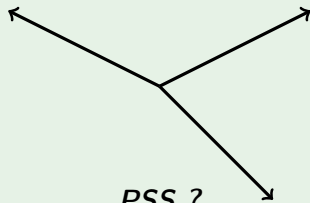
Characterization

M is a PSS \iff There is no $y \neq 0$ such that $y^T M \geq 0^T$.

Examples in \mathbb{R}^2



PSS ?



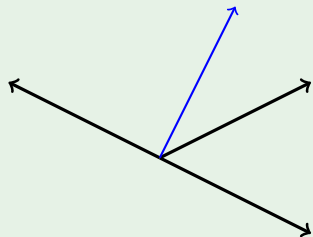
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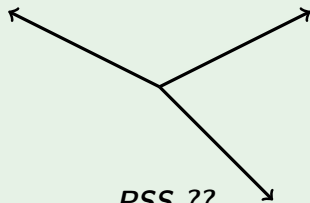
Characterization

M is a PSS \iff There is no $y \neq 0$ having an acute angle with every element of M .

Examples in \mathbb{R}^2



not a PSS !



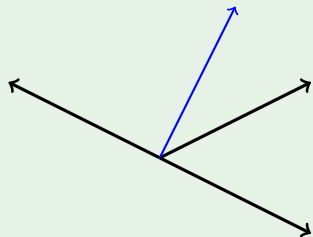
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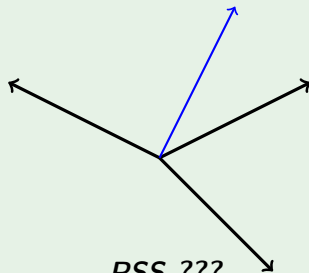
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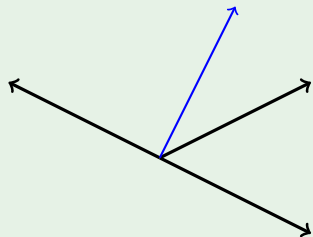
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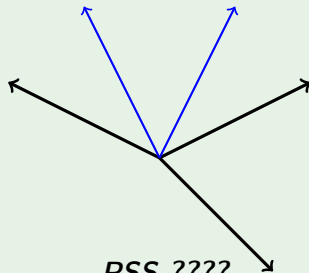
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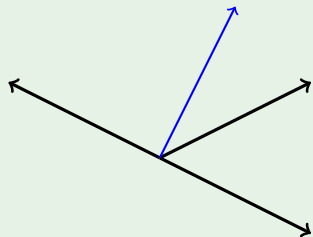
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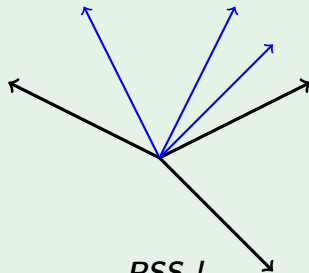
Characterization

M is a PSS \iff There is no $y \neq 0$ having an acute angle with every element of M . *Only a finite number of checks are required !*

Examples in \mathbb{R}^2



not a PSS !



PSS !

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Characterizing PSSs via the cosine measure

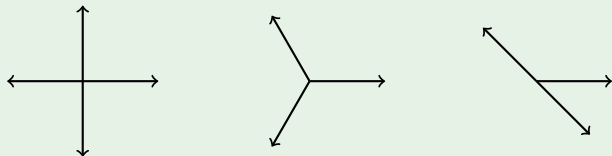
Definition

Let $M = \{d_1, \dots, d_m\} \subset \mathbb{R}^n$. The **cosine measure** of M is defined as:

$$cm(M) := \min_{v \neq 0} \max_{i \in [1, m]} \frac{d_i^\top v}{\|d_i\| \cdot \|v\|}.$$

Characterization

M is a PSS $\iff cm(M) > 0$.



Characterizing PSSs via the cosine measure

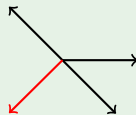
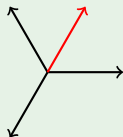
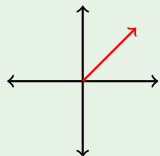
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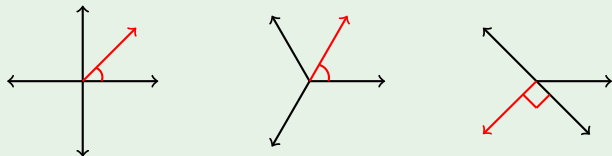
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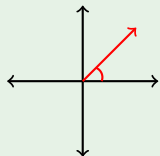
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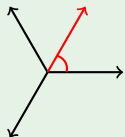
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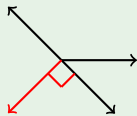
M is a PSS $\iff cm(M) > 0$.



$$cm(M) = \frac{1}{\sqrt{2}}$$



$$cm(M) = \frac{1}{2}$$



$$cm(M) = 0$$

Computing the cosine measure

Theorem

Let $D_{\mathbb{R}^n}$ be a positive basis of \mathbb{R}^n .

- If $|D_{\mathbb{R}^n}| = 2n$, then $cm(D_{\mathbb{R}^n}) \leq \frac{1}{\sqrt{n}}$.
- If $|D_{\mathbb{R}^n}| = n + 1$ then $cm(D_{\mathbb{R}^n}) \leq \frac{1}{n}$.

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- It is better to use PSSs whose cosine measure is close to 1 for optimization purposes...
 - ...However, algorithms to find the cosine measure of a PSS are exponential in time.

Orthogonally structured positive bases

Definition

Positive basis $D_{\mathbb{R}^n}$ of \mathbb{R}^n is an OSPB if:

- It can be written as a partition of positive bases for smaller linear spaces:

$$D_{\mathbb{R}^n} = D_{\mathbb{L}_1} \cup \dots \cup D_{\mathbb{L}_s}.$$

- These bases are pairwise orthogonal and of minimal size.

Examples

$$\begin{bmatrix} 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{bmatrix} \quad \begin{bmatrix} 0 & 1 & 4 & 0 & -5 \\ 2 & 0 & 0 & -2 & 0 \\ 0 & 3 & 5 & 0 & -8 \end{bmatrix} \quad \begin{bmatrix} 1 & -1 & 2 & -1 & -1 \\ 1 & -1 & -1 & 1 & 0 \\ 1 & -1 & -1 & 0 & 1 \end{bmatrix}.$$

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Theorem (new !)

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Cosine measure of an OSPB

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- *Step 2: For any set $D_{\mathbb{L}_i}$ in the decomposition, compute its induced cosine measure c_i .*
- *Step 3: Return $\frac{1}{\sqrt{\sum_i c_i^{-2}}}$.*

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What if we have trouble computing $f(x_0 + \alpha d_i)$?

Positive k -spanning sets, Positive k -bases

Definition

A **PkSS** remains positively spanning when $k - 1$ of its elements are removed. A **positive k -basis** is a minimal PkSS.

Examples of positive 2-bases in \mathbb{R}^2

$$\begin{bmatrix} 1 & 0 & -1 & 1 & 0 & -1 \\ 0 & 1 & -1 & 0 & 1 & -1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & -1 & 0 & 1 & -1 \\ 0 & 1 & 0 & -1 & 1 & -1 \end{bmatrix}$$

Remark

$P1SS \iff PSS$

Positive k -spanning sets, Positive k -bases

Definition

At least k elements must be removed from a **PkSS** before it stops being positively spanning. A **positive k -basis** is a minimal PkSS.

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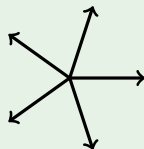
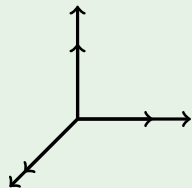
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Alternative definition

Characterization

M is a $PkSS$ \iff For all $y \neq 0$, vector $y^T M$ has at least k positive coordinates.

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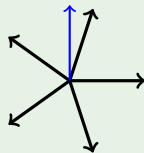
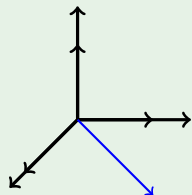


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Examples in \mathbb{R}^2

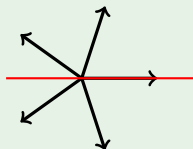
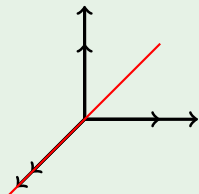


Alternative definition

Characterization

M is a P_k SS \iff For any *hyperplane*, at least k elements of M point on each side of the hyperplane.

Examples in \mathbb{R}^2



Cardinal of positive k -bases

Proposition

Let $D_{\mathbb{R}^n}$ be a positive basis of \mathbb{R}^n . Then $n + 1 \leq |D_{\mathbb{R}^n}| \leq 2n$.

Can this be generalized ?

Let us try...

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- *The hyperplane characterization implies a lower bound of $2k + n - 1$ on the size of a $PkSS$. Tight (Marcus, 1984).*
- *Digraphs can be used to create positive k -bases. In that case, the maximal size is $2kn$.*

PkSSs and Polytope theory

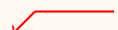
Definition

Let $M \in \mathbb{M}_{n,m}(\mathbb{R})$. Any full rank matrix B such that $MB^T = 0$ is called a **Gale diagram** of M .

Theorem

Suppose $M = \begin{bmatrix} P & & \\ 1 & \dots & 1 \end{bmatrix}$. Then any Gale diagram of M is a PkSS.

Vertices of a well chosen polytope.



Polytopes and cardinality of positive k -bases

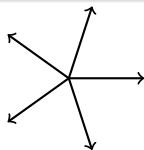
Theorem

Let $D_{\mathbb{R}^n}^{(k)}$ be a positive k -basis of \mathbb{R}^n . Then :

$$2k + n - 1 \leq |D_{\mathbb{R}^n}^{(k)}| \leq kn(n + 1)^{k-1}$$

Remark

- The lower bound is tight (eg: pentagon).
- The upper bound might not be. (Wotzlaw, 2009)



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A new tool for characterizing PkSSs

Definition (new !)

Let $k \geq 1$ and $M \subset \mathbb{R}^n \setminus \{0\}$. The k -cosine measure of M is:

$$cm_k(M) := \min_{v \neq 0} \max_{\substack{S \subset M \\ |S|=k}} \min_{s \in S} \frac{s^\top v}{\|s\| \cdot \|v\|}.$$

A new tool for characterizing PkSSs

Definition (new !)

Let $k \geq 1$ and $M \subset \mathbb{R}^n \setminus \{0\}$. The k -cosine measure of M is:

$$cm_k(M) := \min_{v \neq 0} \max_{\substack{S \subset M \\ |S|=k}} \min_{s \in S} \frac{s^\top v}{\|s\| \cdot \|v\|}.$$

Remark

$$cm_1(M) = cm(M).$$

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Theorem

M is a PkSS if and only if $cm_k(M) > 0$.

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$cm_1(M) = cm(M)$.

Alternative definition

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Let $k \geq 1$ and $M \subset \mathbb{R}^n \setminus \{0\}$. Then:

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- $cm_k(M)$ is computed by finding a vector as far as possible from its k^{th} closest neighbor in M .

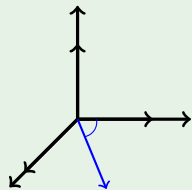
Examples

Remark

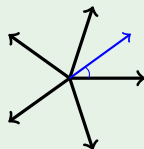
$cm_1(M)$ is computed by finding a vector as far as possible from its closest neighbor in M .

Computing the cosine measure:

Examples



$$\cos\left(\frac{3\pi}{8}\right) \simeq 0.38$$



$$\cos\left(\frac{2\pi}{10}\right) \simeq 0.809$$

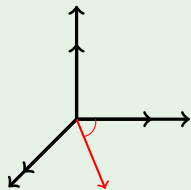
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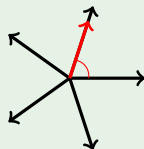
$cm_2(M)$ is computed by finding a vector as far as possible from its *second* closest neighbor in M .

Computing the 2-cosine measure:

Examples



$$\cos\left(\frac{3\pi}{8}\right) \simeq 0.38$$



$$\cos\left(\frac{2\pi}{5}\right) \simeq 0.309$$

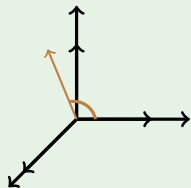
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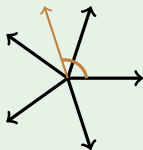
$cm_3(M)$ is computed by finding a vector as far as possible from its *third* closest neighbor in M .

Computing the 3-cosine measure:

Examples



$$\cos\left(\frac{5\pi}{8}\right) \simeq -0.38$$



$$\cos\left(\frac{3\pi}{5}\right) \simeq -0.309$$

Basic remarks of the k -cosine measure

There is no easy way to compute the k -cosine measure of a given family.
Let us focus on finding bounds.

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- *Let $k < l$. Then $cm_k(M) \geq cm_l(M)$.*
- *Let $N \subseteq M$. Then $cm_k(N) \leq cm_k(M)$.*
- *If $cm(M) = \alpha$, duplicating M creates a PkSS whose k -cosine measure is α .*

Rotating PSSs

Proposition

Let $D_{\mathbb{R}^n} \subset \mathbb{R}^n \setminus \{0\}$ be an OSPB. Stacking together k rotations of $D_{\mathbb{R}^n}$ creates a PkSS $D_{\mathbb{R}^n}^{(k)}$ satisfying:

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Remark

Unfortunately, this PkSS might not be a positive k -basis.

Rotating minimal bases

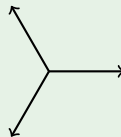
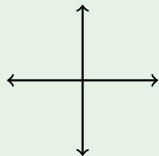
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Rotating minimal positive bases can create positive k -bases.

Idea: apply small rotations.

Note: this technique only works for minimal positive bases !

Examples in \mathbb{R}^2



Rotating minimal bases

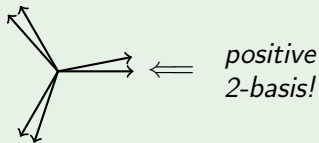
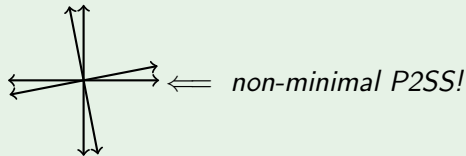
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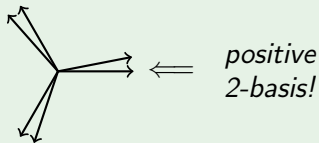
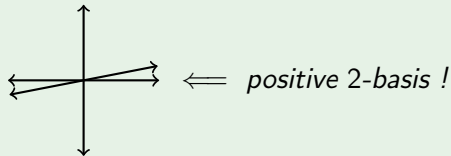
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- Finding new ways to build "nice" PkSSs.

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Perspectives

- Finding new ways to build "nice" PkSSs.
- Using PkSSs in Derivative Free Algorithms.

Thanks for your attention !